

Let's imagine that we introduce a new coin system. Instead of using pennies, nickels, dimes, and quarters, let's say we agree on using only 4-cent and 7-cent coins. One might point out the following flaw of this new system: certain amounts cannot be exchanged, for example, 1, 2, or 5 cents. On the other hand, this deficiency makes our new coin system more interesting than the old one, because we can ask the question: "which amounts can be changed?" We will see shortly that there are only finitely many integer amounts that *cannot* be exchanged using our new coin system. A natural question, first tackled by Ferdinand Georg Frobenius and James Joseph Sylvester in the nineteenth century, is: "what is the *largest* amount that cannot be exchanged?" As mathematicians, we like to keep questions as general as possible, and so we ask: given coins of denominations a and b—positive integers without a common factor—can you give a formula g(a, b) for the largest amount that cannot be exchanged using the coins  $a_1, a_2, \ldots, a_n$  is known as the *Frobenius coin-exchange problem*. To study the Frobenius number g(a, b), we use the *Euclidean Algorithm*. For integers a and b that have no common factor, this algorithm yields integers x and y such that ax + by = 1.

## Problems

1. Find g(4,7) and g(5,11).

2. Find x and y such that 4x + 7y = 1. Find another x and y such that 4x + 7y = 1.

3. Find x and y such that 5x + 11y = 1. Find x and y such that 5x + 11y = 39.

4. Show that, if t is a given integer, we can always find integers x and y such that 4x + 7y = t. Generalize to two coins a and b with no common factor.

5. Show that, if t is a given integer, we can always find integers x and y such that 4x + 7y = t and  $0 \le x \le 6$ . Generalize to two coins a and b with no common factor.

6. Show that the following recipe for determining whether or not a given amount t can be changed (using the coins 4 and 7) works: Given t, find integers x and y such that 4x + 7y = t and  $0 \le x \le 6$ . Then t can be changed precisely if  $y \ge 0$ . Generalize to two coins a and b with no common factor.

7. Use the previous argument to re-compute g(4,7). Generalize your argument to compute g(a,b), for any two coins a and b with no common factor.

8. Suppose t is an integer between 1 and ab - 1 that is not a multiple of a or b. Prove that if the amount t can be changed then ab - t cannot be changed, and conversely, if t cannot be changed then ab - t can be changed.

9. Prove that there are  $\frac{1}{2}(a-1)(b-1)$  amounts that cannot be changed.

10. Think about why g(a, b) actually exists, if a and b have no common factor. More generally, prove that the general Frobenius problem is well defined. That is, show that, given  $a_1, a_2, \ldots, a_d$  with no common factor, every sufficiently large integer is representable (in terms of  $a_1, a_2, \ldots, a_d$ ).

11. Next week we will study the counting sequence

$$r_k = \# \{ (m, n) \in \mathbb{Z}^2 : m, n \ge 0, ma + nb = k \}$$

In words,  $r_k$  counts the representations of  $k \in \mathbb{Z}_{\geq 0}$  as nonnegative linear combinations of a and b. The Frobenius problem asks for the largest among the  $r_k$ 's that is 0. Prove that  $r_{k+ab} = r_k + 1$ .

## A few remarks

The simple-looking formula for g(a, b) that you have found in () inspired a great deal of research into formulas for the general Frobenius number  $g(a_1, a_2, \ldots, a_d)$ , with limited success: While it is safe to assume that the formula for g(a, b) has been known for more than a century, no analogous formula exists for  $d \ge 3$ . The case d = 3 is solved algorithmically, i.e., there are efficient algorithms to compute g(a, b, c) [2, 4, 5], and in form of a semi-explicit formula [3, 7]. The Frobenius problem for fixed  $d \ge 4$  has been proved to be computationally feasible [1, 6], but not even an efficient practical algorithm for d = 4 is known. The formula in () is due to Sylvester and was published as a math problem in the *Educational Times* more than a century ago [9]. For more on the Frobenius problem, we refer to the research monograph [8]; it includes more than 400 references to articles written about the Frobenius problem.

## References

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